

# 图数据库 Neo4j 的实践之路

魏佳  
技术主管 @ LinkedIn中国

# 关于我



2009



2012



2014



2015

IBM  
资深工程师

新浪微博  
技术专家

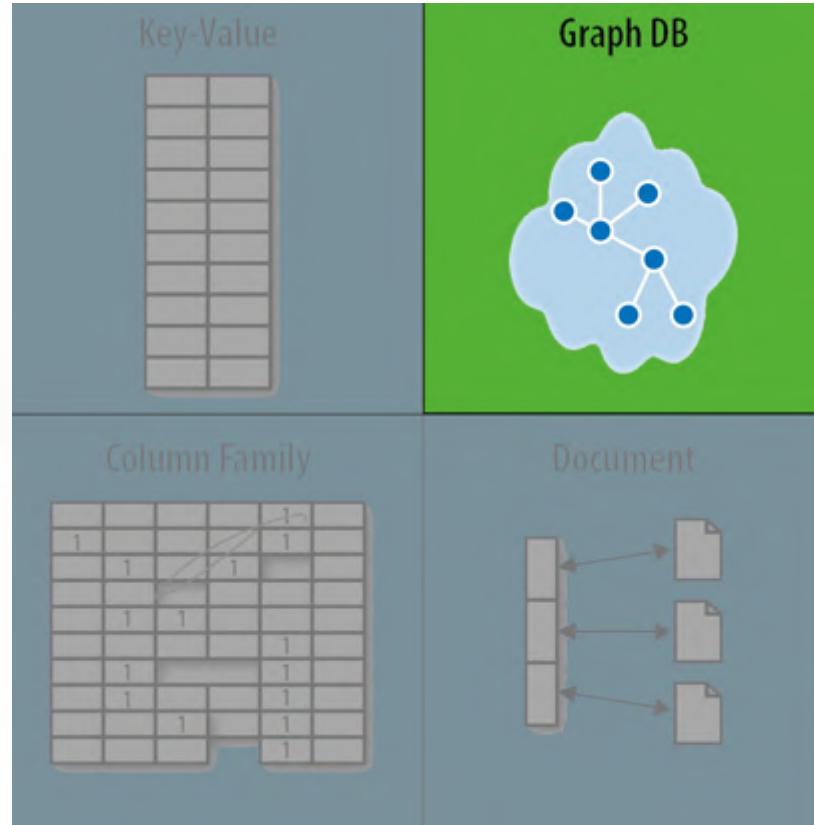
Mico  
联合创始人/CTO

LinkedIn  
技术主管

- 应用生命周期管理 (ALM)
- 分布式消息中间件
- 虚拟化及多租户运行时 (Multi-tenant runtime)
- 分布式系统请求链路追踪和服务质量 (SLA)
- 著有多篇国内外技术专利

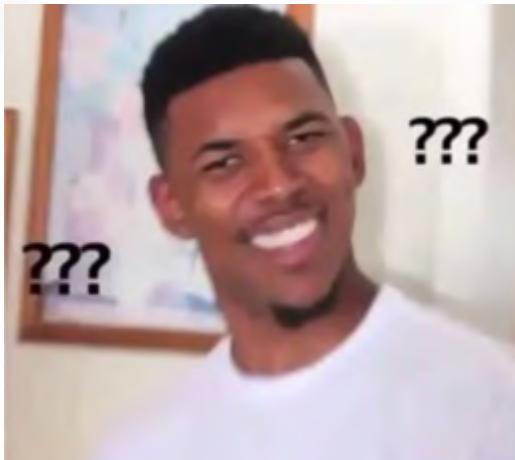
# 关于此分享

- 内容
- 受众



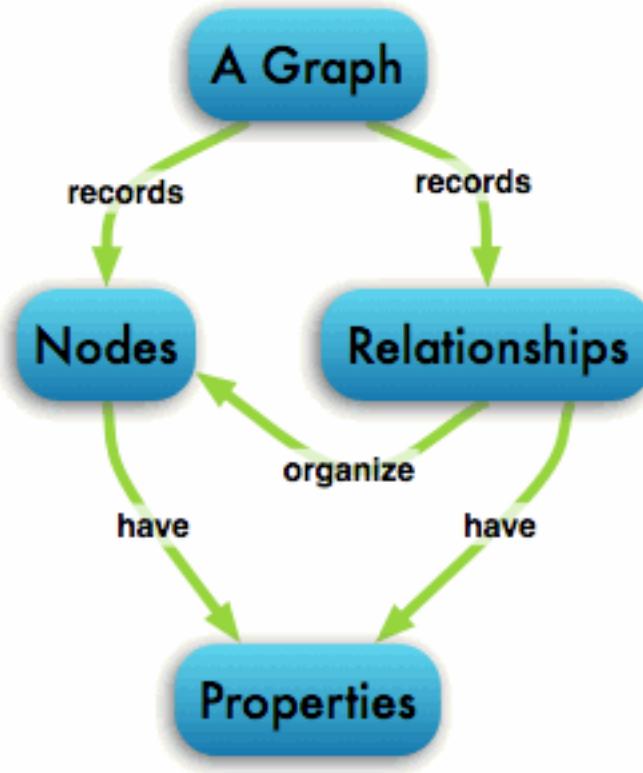
# 图数据库

“...使用图结构 (graph structure) 来实现语义查询，使用点 (vertex) 、边 (edge) 和属性 (property) 来表示和存储数据。最重要的概念是边 (关系)，在存储中，它直接将数据项关联起来...”



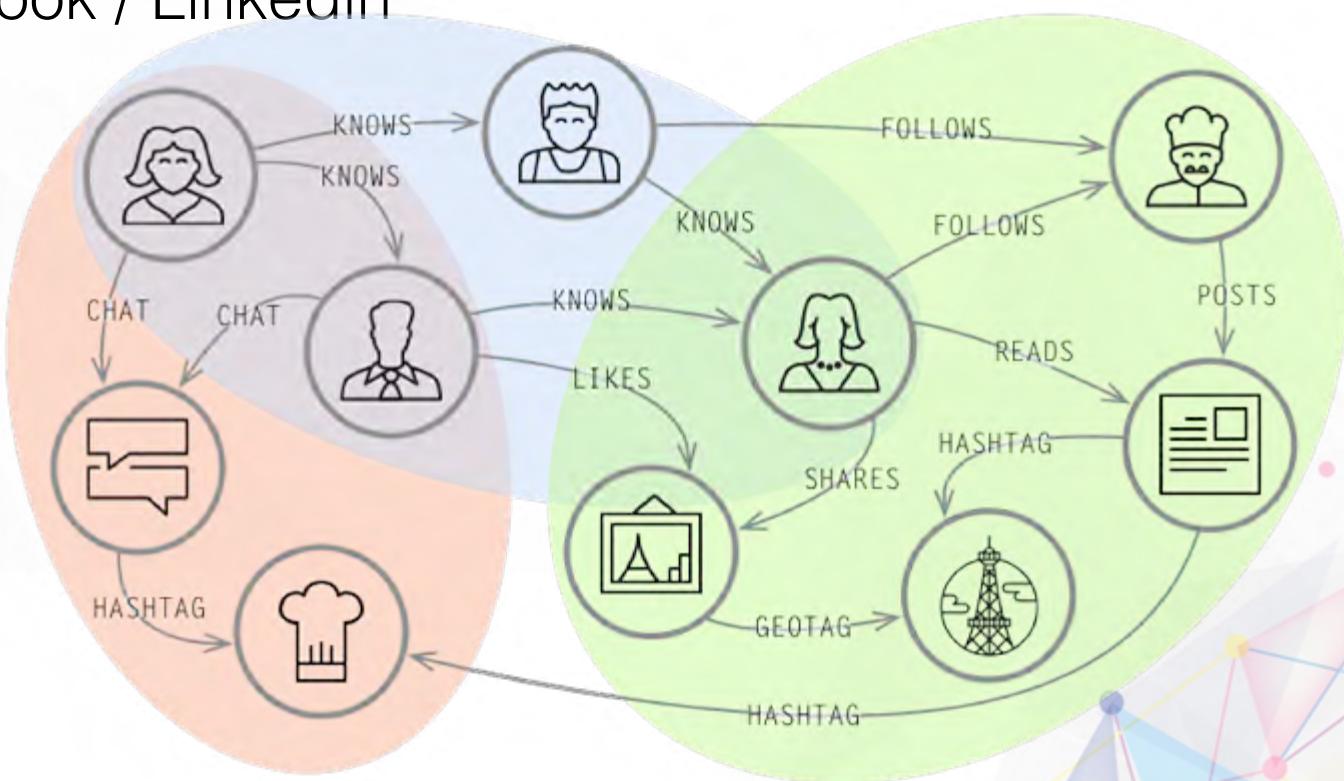
[https://en.wikipedia.org/wiki/Graph\\_database](https://en.wikipedia.org/wiki/Graph_database)

# 图数据库



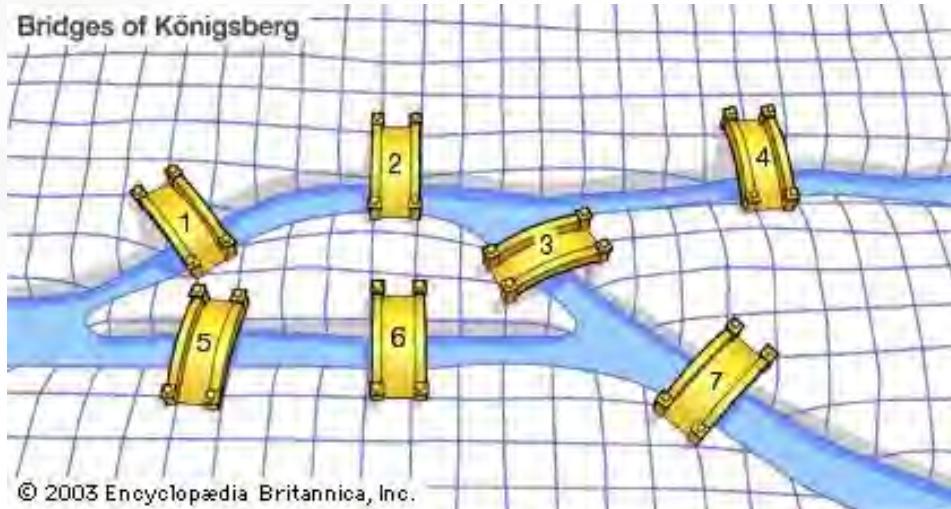
# 场景

- Social Network
  - Facebook / Linkedin



# 场景

- Spatial/GIS
  - 路径规划
  - 交通预测



## Problem Definition

- The traveling salesman problem consists of a salesman and a set of cities. The salesman has to visit each one of the cities starting from a certain one (e.g. the hometown) and returning to the same city. The challenge of the problem is that the traveling salesman wants to minimize the total length of the trip.

Example:

Consider the following set of cities:

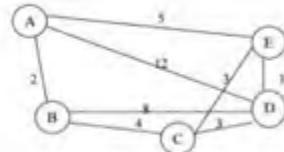


Figure 10.1 A graph with weights on its edges

The problem lies in finding a minimal path passing from all vertices once. For example the path Path1 {A, B, C, D, E, A} and the path Path2 {A, B, C, E, D, A} pass all the vertices but Path1 has a total length of 24 and Path2 has a total length of 31.

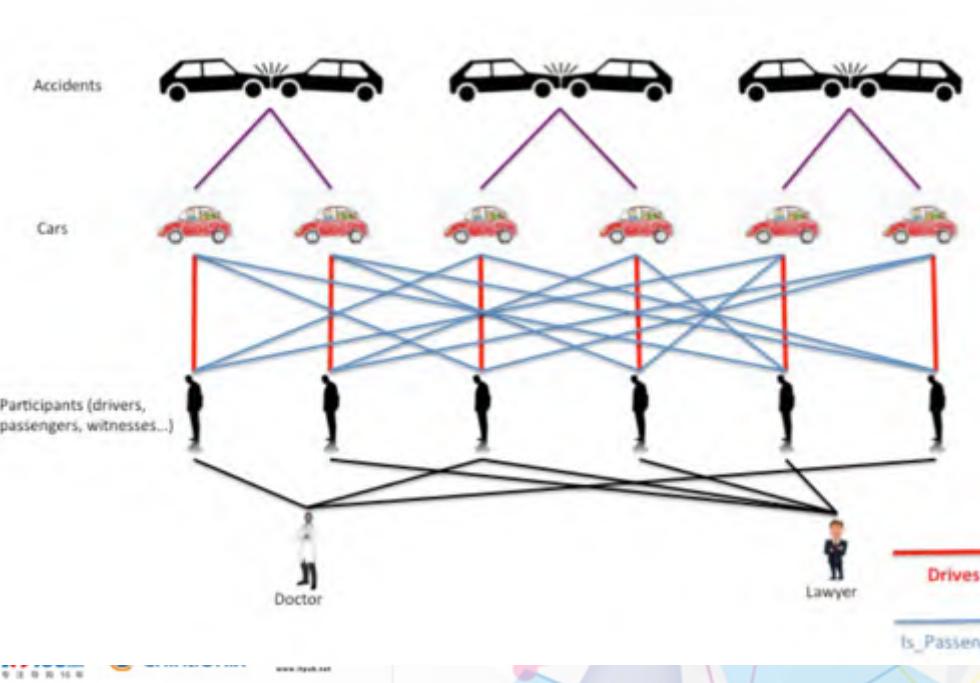
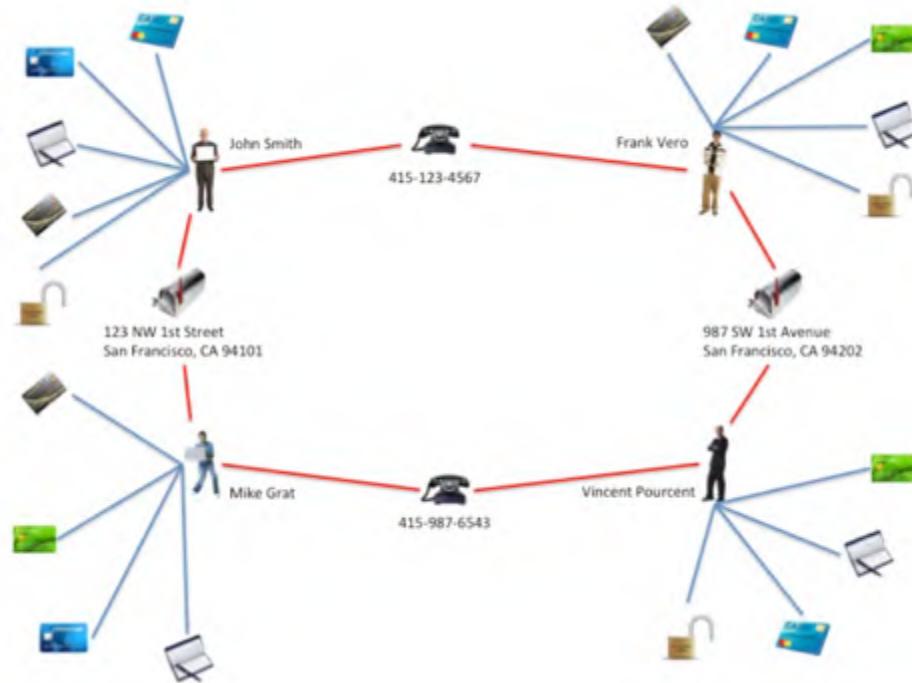
# 场景

- Recommendation
  - “You may also like...”
  - “People you may know...”



# 场景

- Fraud detection
  - First-party bank fraud/Insurance fraud/e-Commerce fraud



# 场景

- Workflow engine
- Access control
- ...

# 场景

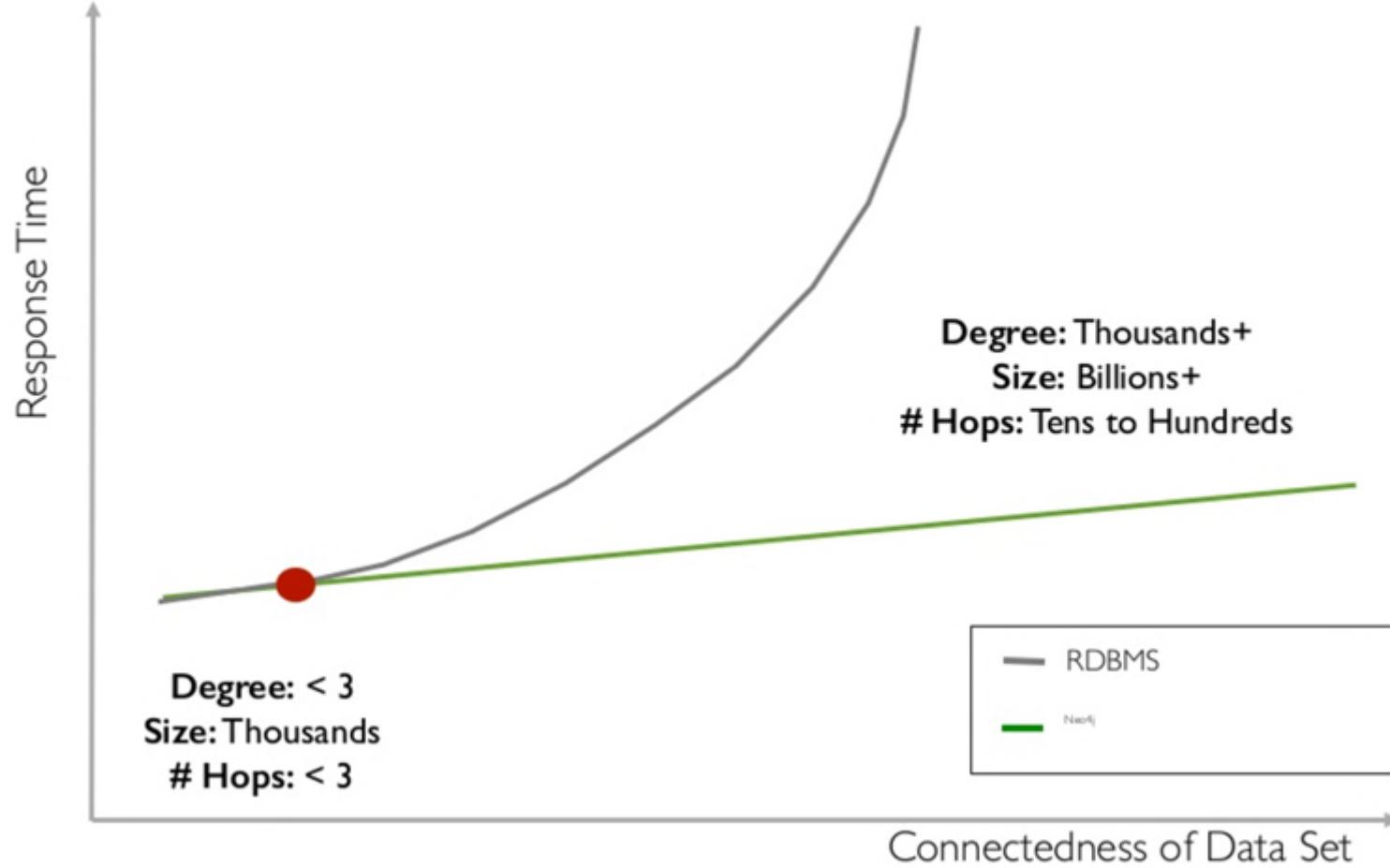


# 场景

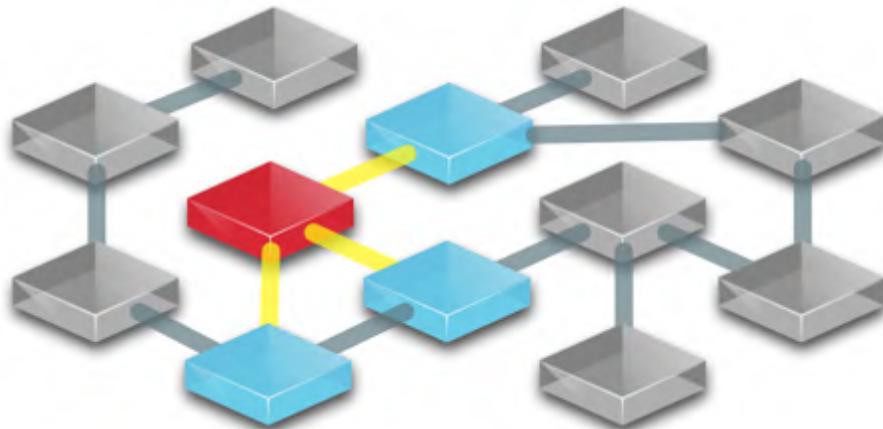
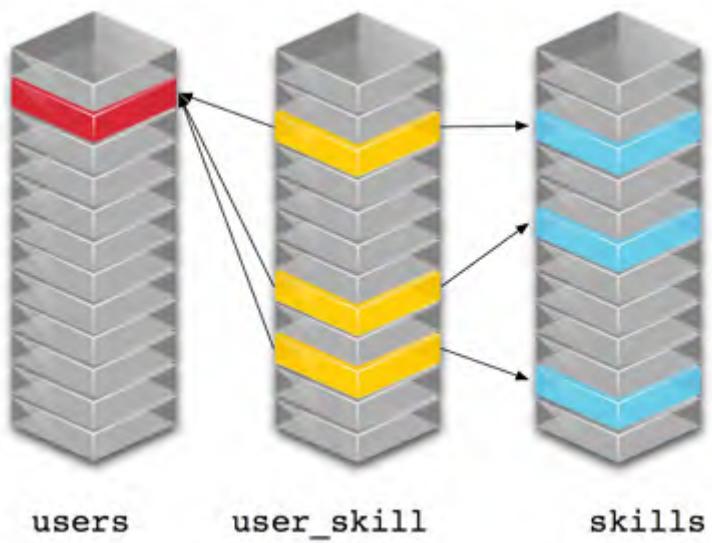
Graph Query Response Time =  
 $f(\text{graph density}, \text{graph size}, \text{query degree})$

- graph density: avg. (relations per node)
- graph size:  $\sum(\text{node})$
- query degree: hops per query

# 场景



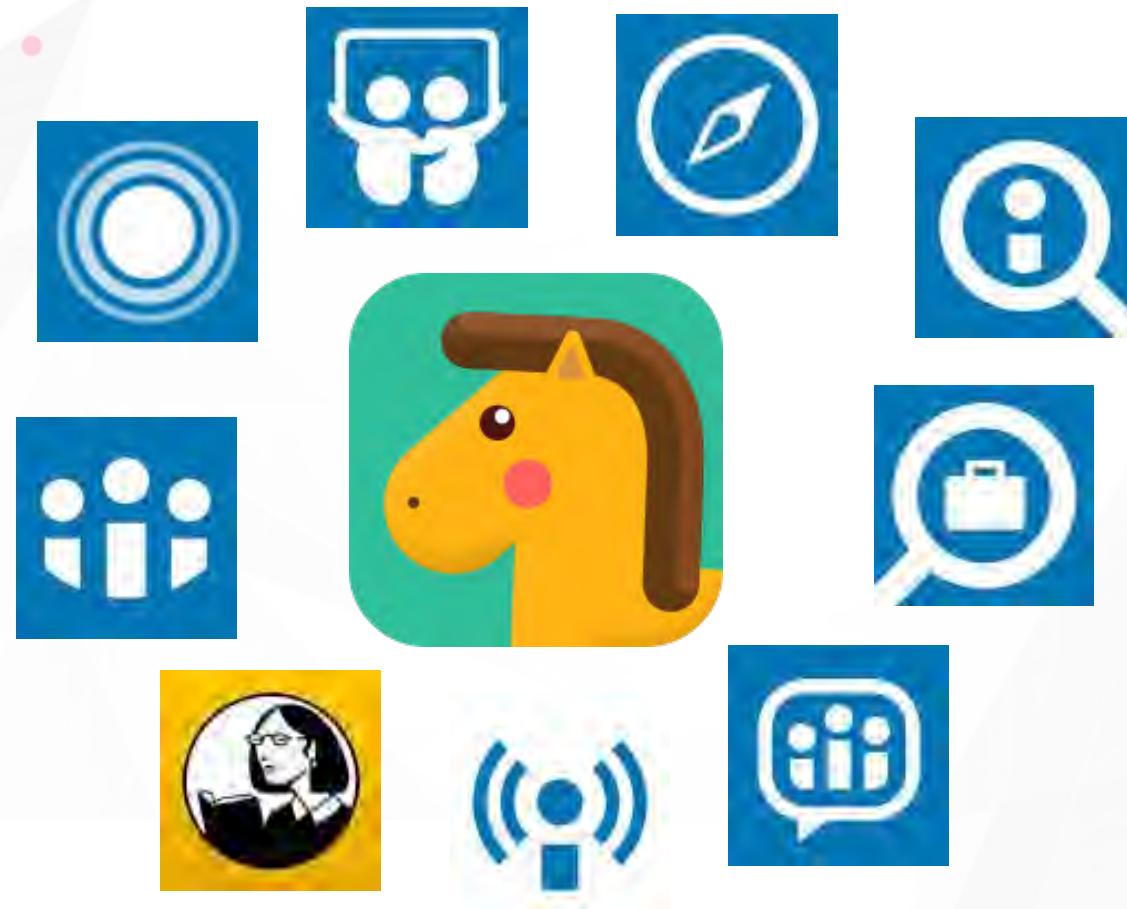
# 场景



# 背景

- LinkedIn Flagship (职场人士)
  - LinkedIn Students (学生群体)
  - LinkedIn Influencer (职场影响力内容发布和传播)
  - LinkedIn Job Search
  - LinkedIn Recruiter (企业征才解决方案)
  - LinkedIn Lookup (企业雇员信息解决方案)
- Lynda (全球最大在线职业教育培训)
- Slideshare (全球最大的在线演示文档共享)
- Pulse (阅读和内容聚合)

# 背景



# 选型



选型



TITAN

# 选型

- OrientDB
  - 缺少批量写，从DB文件加载并初始化非常慢
  - 在200M+节点后写瓶颈
  - 不健壮不稳定
  - CEO不靠谱 <http://orientdbileaks.blogspot.com/2015/06/the-orientdb-issues-that-made-us-give-up.html>



# 选型

- Titan
  - 更像是API layer，可选backend engine，Cassandra、HBase、BDB和ElasticSearch
  - ACID不完备，依赖backend engine
  - Key index使用限制较多
  - 缺乏客户案例
  - Dev freeze



# 选型

The last man standing, and the  
winner is



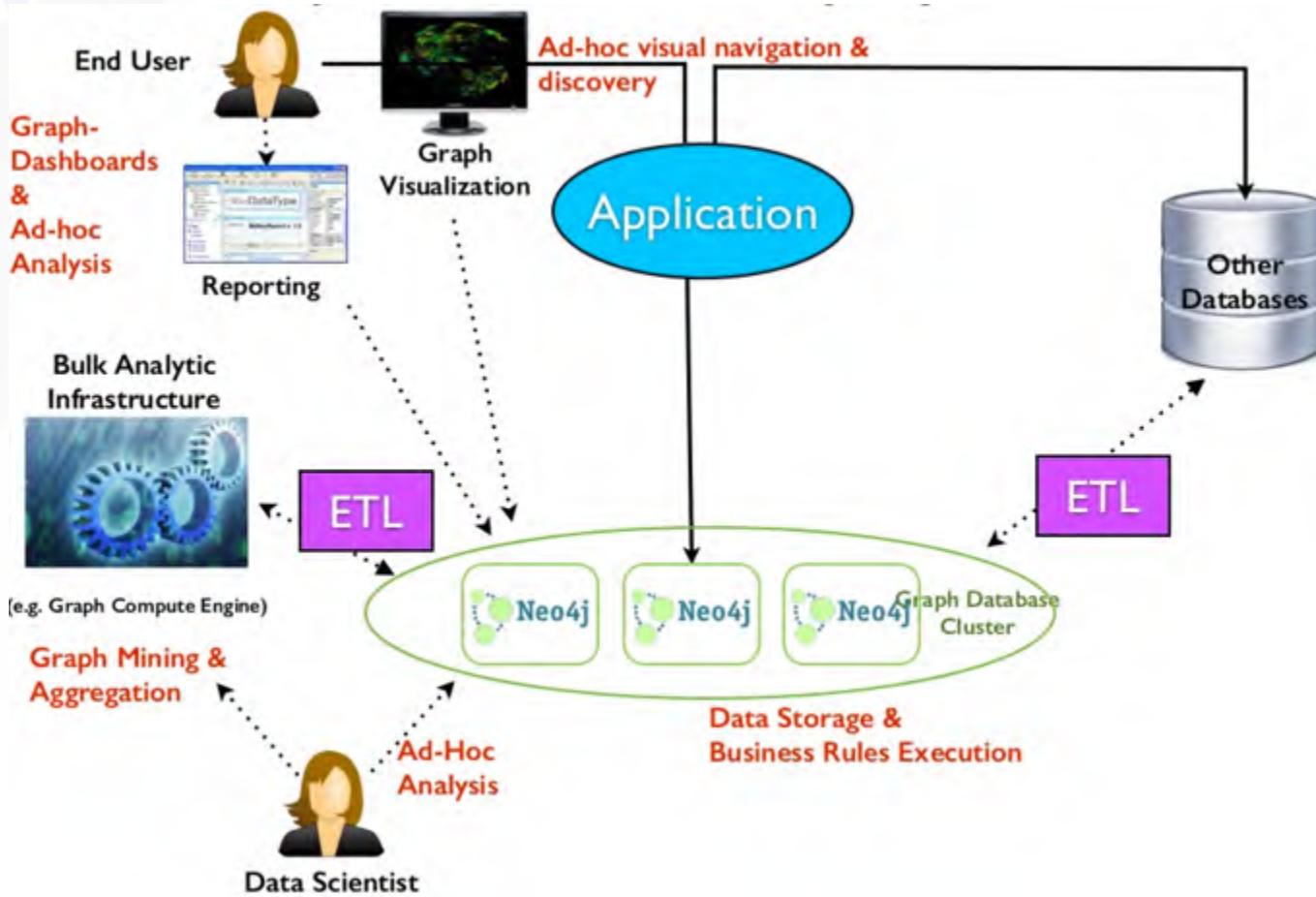
# Neo4j



# Neo4j

- Schema-free with labeled property
- Fast graph traversal(~2M/s)
- Written in Java
- Full ACID compliant
- Full/incremental online backup
- HTTP API or embedded in JVM
- Powerful declarative query language
- Cache(chunk) based sharding
- Linearly scalable read, M-S replication, fault-tolerance
- Plugins & extensions
- ...

# Neo4j



# Neo4j

- Distributed
  - replication
  - fault tolerance
  - scalability
  - consistency

# Neo4j

- Cluster Management
  - member join/leave/heartbeat
- Failover
  - election
  - status propagation
- Replication
  - distributed lock
  - transaction pull/push



Apache  
Zookeeper

HELIX

# Neo4j

- Multi-Paxos algorithm
  - less complex and testable
- Ready for cluster Management and election

# Neo4j

```
public enum AcceptorState
    implements State<AcceptorContext, AcceptorMessage> {
start {
    public AcceptorState handle(AcceptorContext context, Message<AcceptorMessage> message, MessageHolder outgoing)
        switch (message.getMessageType()) {
            case join: {
                return acceptor; ← new state
            }
        }
    return this; ← no state change for fallthrough
},
acceptor { ← self == current state
    public AcceptorState handle(AcceptorContext context, Message<AcceptorMessage> message, MessageHolder outgoing)
        switch (message.getMessageType()) {
            case prepare: {...}
            case accept: {...}
            case leave: {
                context.leave(); ← actions
                return start;
            }
        }
    return this;
}
```

Annotations:

- new state**: Points to the line `return acceptor;`
- no state change for fallthrough**: Points to the line `return this;`
- self == current state**: Points to the line `acceptor {`
- mutable data**: Points to the line `switch (message.getMessageType()) {`
- message w/ type & payload**: Points to the line `case prepare: {...}`
- actions**: Points to the line `context.leave();`
- receiver of outgoing messages**: Points to the line `MessageHolder outgoing`

# Neo4j

```
public enum HeartbeatState
    implements State<HeartbeatContext, HeartbeatMessage> { state machine
start {...},
state
heartbeat {
    public HeartbeatState handle(HeartbeatContext context, Message<HeartbeatMessage> message, MessageHolder outgoing)
        final ClusterContext clusterContext = context.getClusterContext();
        final LearnerContext learnerContext = context.getLearnerContext();
        switch (message.getMessageType()) {
            case i_am_alive: {...}
            case timed_out: {...}
            case sendHeartbeat: {
                URI to = message.getPayload();

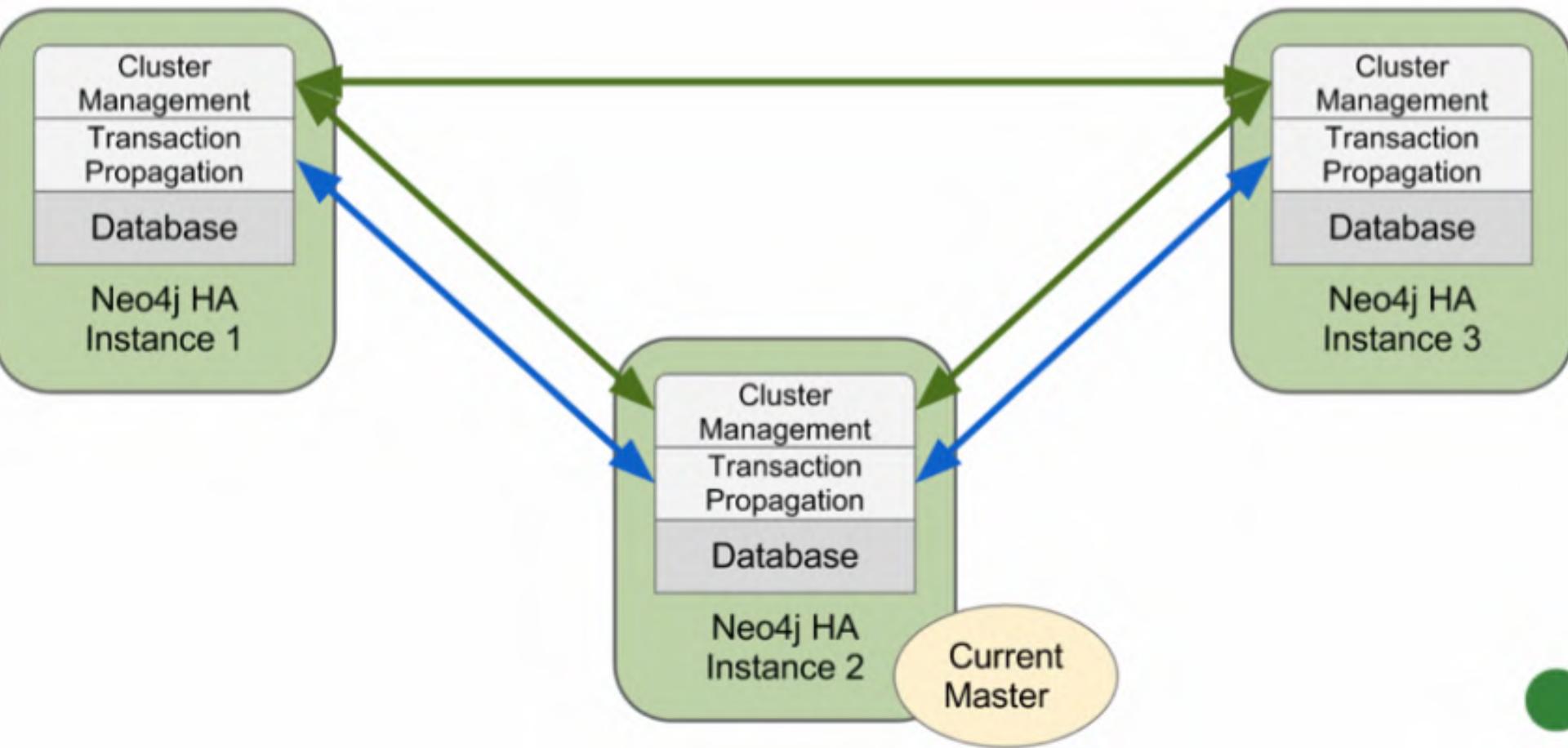
                // Check if this node is no longer a part of the cluster
                if (clusterContext.getConfiguration().getMembers().contains(to)) {
                    // Send heartbeat message to given server
                    outgoing.offer(to(HeartbeatMessage.i_am_alive, to,
                        new HeartbeatMessage.IAmAliveState(clusterContext.getMe())))
                        .setHeader("last-learned", learnerContext.getLastLearnedInstanceId() + "");

                    // Set new timeout to send heartbeat to this host
                    clusterContext.timeouts.setTimeout(HeartbeatMessage.sendHeartbeat + "-"
                        + to, timeout(HeartbeatMessage.sendHeartbeat, message, to));
                }
                break;
            }
            case reset_send_heartbeat: {...}
            case suspicions: {...}
            case leave: { return start; }
        }
    return this;
}
```

Annotations:

- state machine
- state
- mutable state
- message-payload
- send message
- message-payload
- timeout with message
- new state

# Neo4j

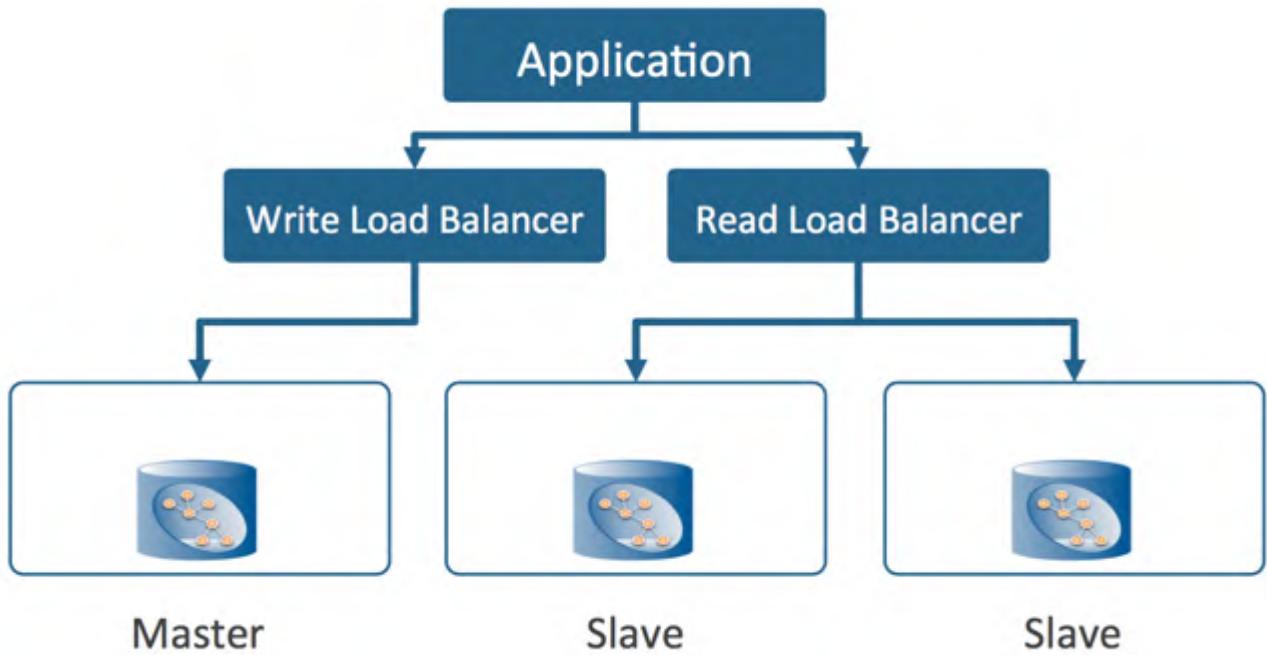


# Neo4j

- All transactions are committed on the master
  - eventually applied to the slaves
  - eventually defined by update interval
- When writing on a slave
  - lock coordinated through the master
  - transaction buffered on the slave
  - applied on the master to obtain txid
  - applied to the slaves

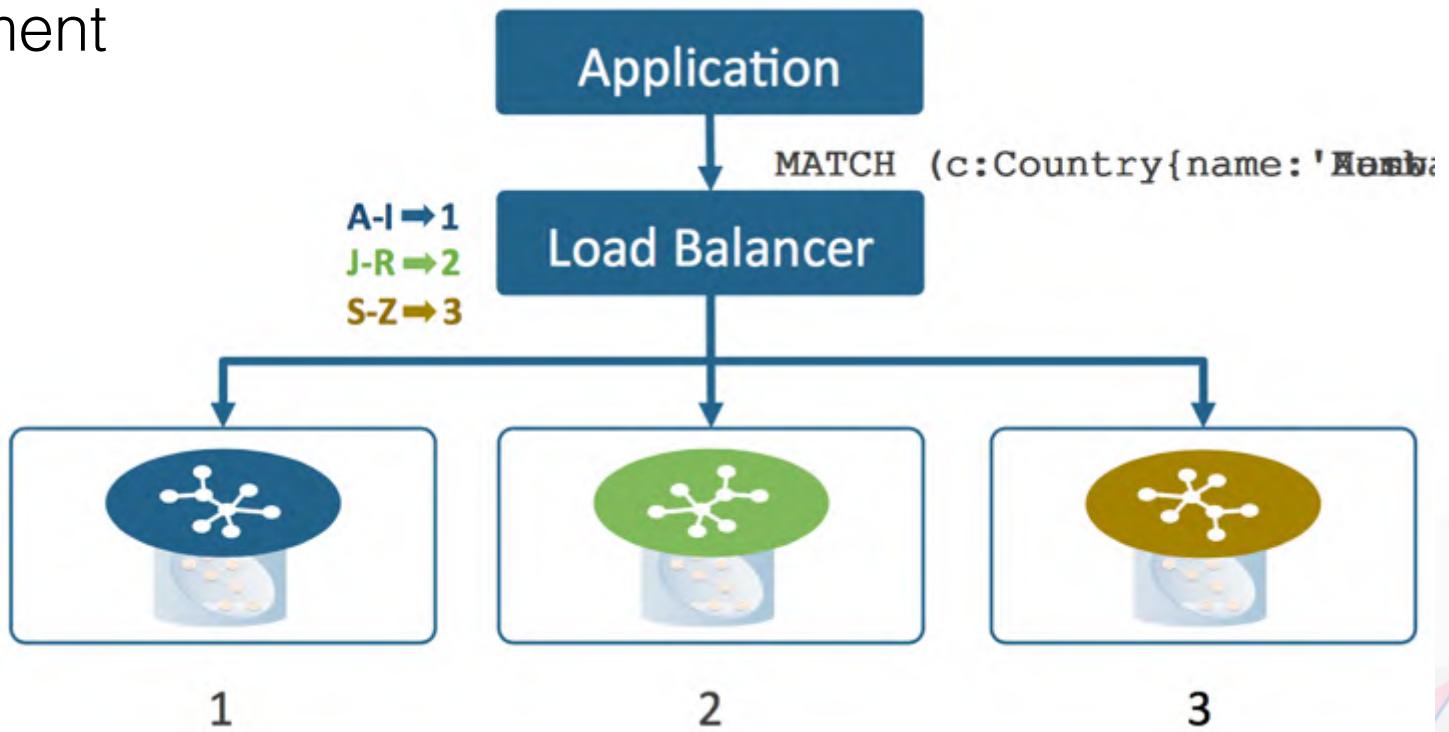
# Neo4j

- Deployment



# Neo4j

- Deployment



# 踩坑



当心坑洞

# 踩坑

- GC暂停导致重选举
  - 进而一致性检查，引起数据复制，最终造成若干节点不可用
- JSON payload大，影响吞吐同时浪费带宽
  - 同时heap中产生大量脏footprint，对GC不利
- HTTP based API效率不高
  - 短链接，协议处理的开销
- 复杂查询场景效率不高
  - 依赖外部其他资源时，结果集利用率低
  - 多次请求时，不必要的网络延时

# 踩坑

- GC暂停导致重选举
  - 进而一致性检查，引起数据复制，最终造成若干节点不可用
- JSON payload大，影响吞吐同时浪费带宽
  - 同时heap中产生大量脏footprint，对GC不利
- HTTP based API效率不高
  - 短链接，协议处理的开销
- 复杂查询场景效率不高
  - 依赖外部其他资源，结果集利用率低
  - 多次请求，不必要的网络延时

# 改造和优化

- GC调优
  - G1收集器
    - adaptive
  - JFR+JMC
    - on-flight
    - visualized
    - insignificant performance overhead

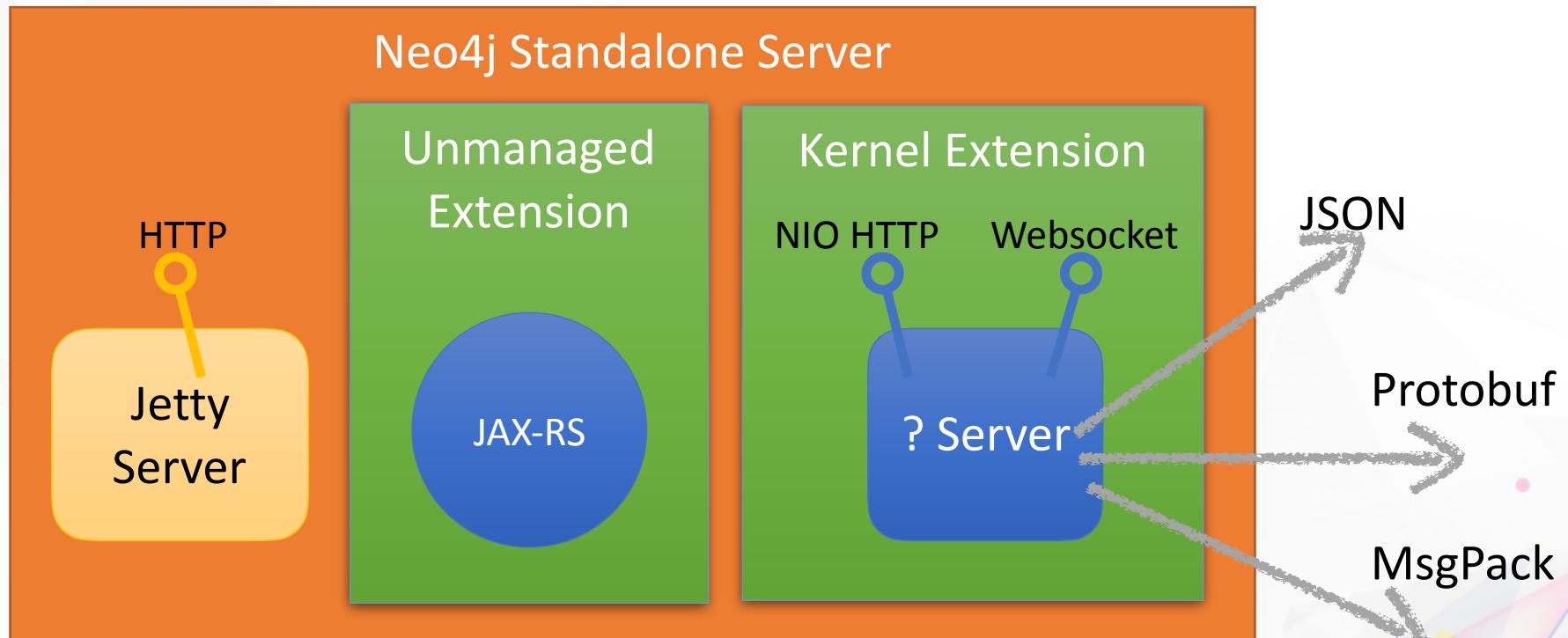
# 踩坑

- GC暂停导致重选举
  - 进而一致性检查，引起数据复制，最终造成若干节点不可用
- JSON payload大，影响吞吐同时浪费带宽
  - 同时heap中产生大量脏footprint，对GC不利
- HTTP based API效率不高
  - 短链接，协议处理的开销
- 复杂查询场景效率不高
  - 依赖外部其他资源，结果集利用率低
  - 多次请求，不必要的网络延时

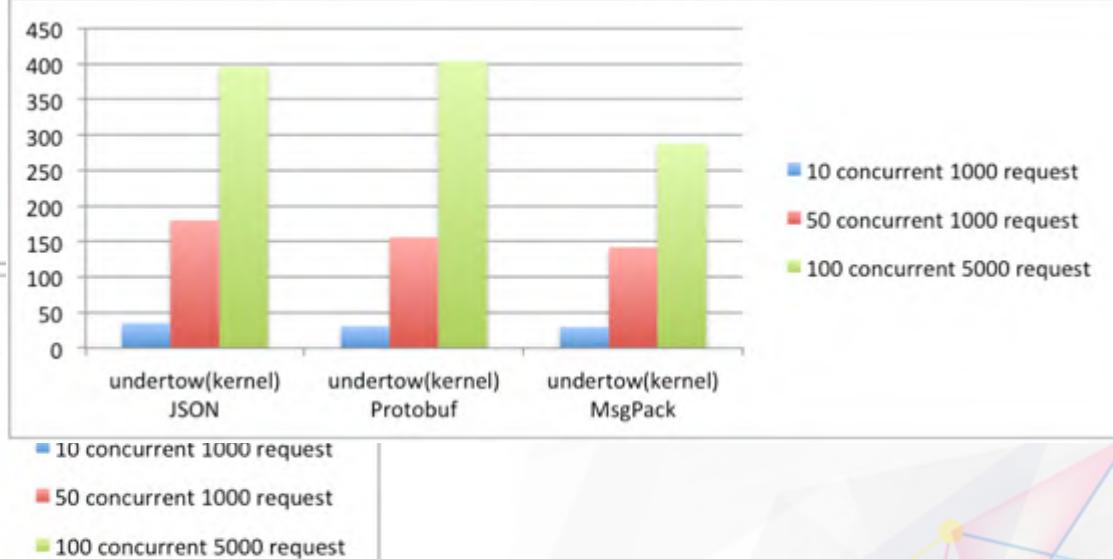
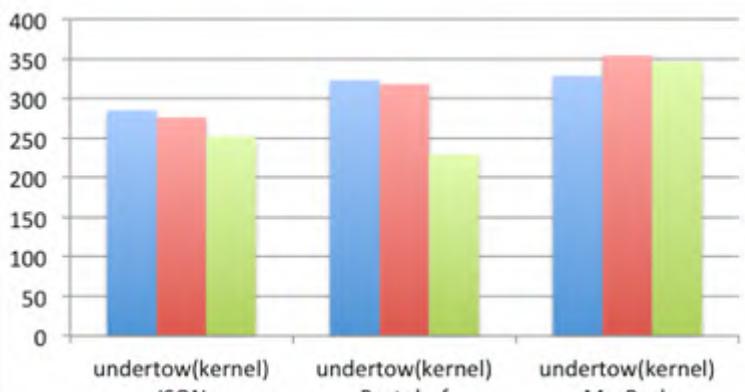
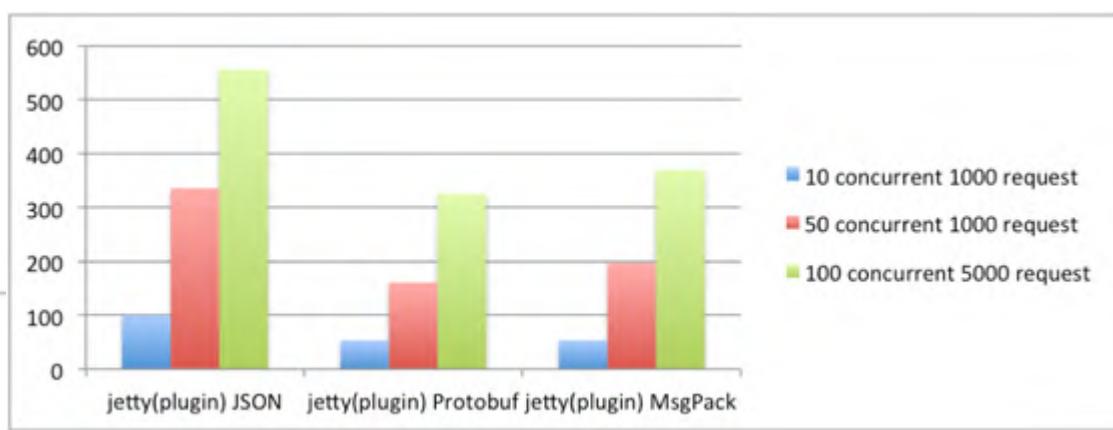
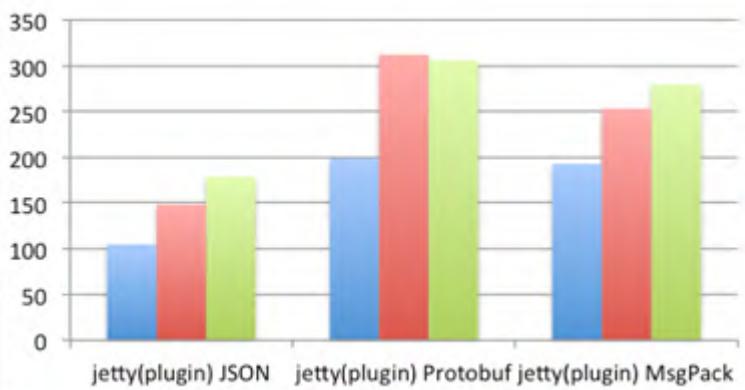
# 改造和优化

- 协议, 传输
  - Jetty+Jersey serve HTTP request/response
  - poor performance
- 数据
  - JSON (jackson streaming)
  - serialization overhead
- Plugin/Extension
  - Plugin
  - Unmanaged extension (JAX-RS)
  - Kernel extension (no constraint)

# 改造和优化

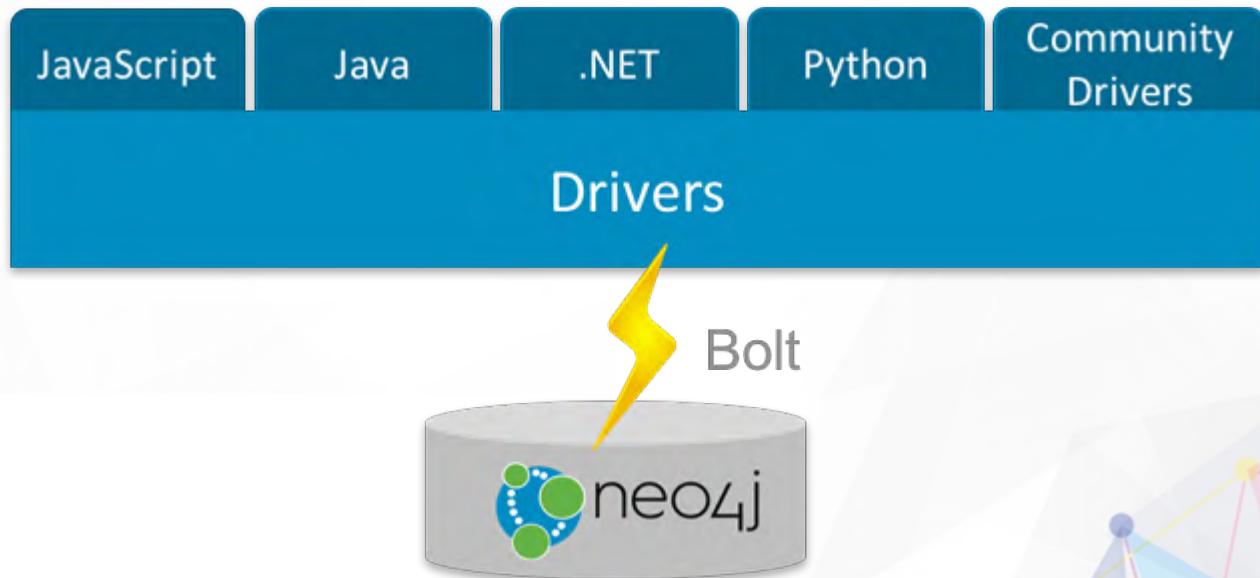


# 改造和优化



# 新特性

- Bolt: Streaming binary protocol
  - TCP persistent oriented
  - PackStream



# 新特性

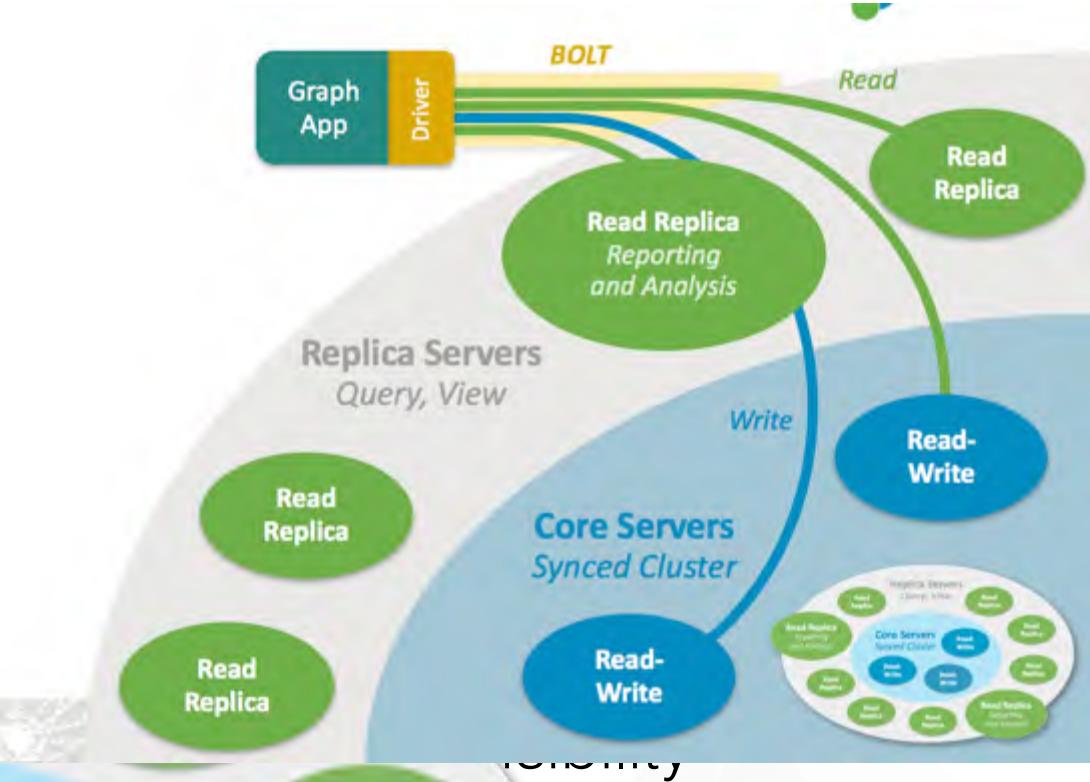
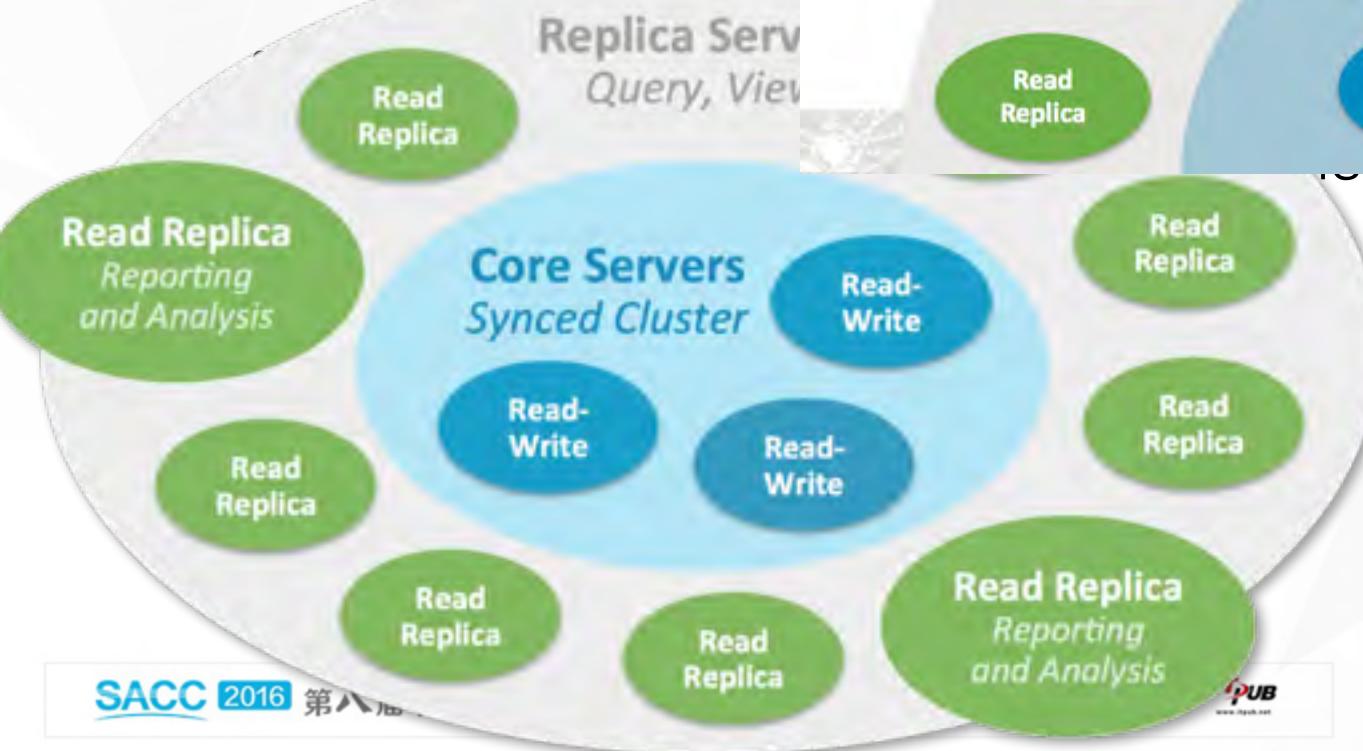
- Paxos->Raft
  - Paxos made *moderately* complex
  - more than “eventual consistency”, causal consistency

# 新特性

- New cluster
  - Core servers提供对等的多主结构，Replica servers只读只参与复制
  - Bolt+driver内置负载均衡的能力
  - Specify guarantees around query visibility
  - Stateful bookmark

# 新特性

- New cluster
    - Core servers 提供  
servers 口述



# 新特性

- Security enhancement
  - role-based authorization
  - support LDAP/Active Directory
- Kernel/storage improvement
  - space compaction
  - automatic execution guard
- Administrative tools
- ...

# 结束

# Q & A



@Life\_Player\_



weijia\_bj

# THANKS

SequeMedia  
盛拓传媒

IT168.com  
专业 专注 16 年

ChinaUnix

ITPUB  
[www.Itpub.net](http://www.Itpub.net)