

### Concurrency in Rust

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#### What's Rust?

Rust is a systems programming language that runs blazingly fast, prevents segfaults, and guarantees thread safety.

#### **Concurrency?**

Rust?

Libraries

Futures

# What's concurrency?

In computer science, concurrency is a property of systems in which several computations are executing simultaneously, and potentially interacting with each other.

# Why concurrency?



# Getting our feet wet

// What does this print?
int main() {
 int pid = fork();
 printf("%d\n", pid);
}

# Concurrency is hard!

- Data Races
- Race Conditions
- Deadlocks

Exploitable!

- Use after free
- Double free 4

#### Concurrency?

#### Rust?

Libraries

Futures

#### What's Rust?

Rust is a systems programming language that runs blazingly fast, **prevents segfaults**, and guarantees thread safety.

# What's **safety**?









```
fn main() {
    let mut v = Vec::new();
    v.push(1);
    v.push(2);
    take(v);
    // ...
fn take(v);
```

}

}



error: use of moved value `v`

# Borrowing

```
fn main() {
    fn pead(v: &Met<\202×132>) {
    let mut v = Vec::new();
    push(&mut v);
    read(&v);
    // ...
}
```



# Safety in Rust

- Rust *statically* prevents aliasing + mutation
- Ownership prevents double-free
- Borrowing prevents use-after-free
- Overall, no segfaults!

#### A data race happens when there are two concurrent memory accesses to the same location in a program where:

Data races

Aliasing!

- at least one is unsynchronized
- at least one is a write Mutation!

#### Concurrency?

Rust?

#### Libraries

Futures

# Rust Concurrency Libs

- Language only provides ownership/borrowing
- Libraries implement common abstractions
- Flexible to cover wide range of paradigms

#### std::thread

#### std::sync::mpsc

let (tx, rx) = mpsc::channel();
let tx2 = tx.clone();
thread::spawn(move || tx.send(5));
thread::spawn(move || tx2.send(4));

// Prints 4 and 5 in an unspecified order
println!("{:?}", rx.recv());
println!("{:?}", rx.recv());

### std::sync::Arc

let shared\_numbers = Arc::new(vec![1, 2, 3]);
let child\_numbers = shared\_numbers.clone();
thread::spawn(move || {
 assert\_eq!(child\_numbers, [1, 2, 3]);
});

assert\_eq!(shared\_numbers, [1, 2, 3]);

#### std::sync::atomic::\*

let number = AtomicUsize::new(10); let prev = number.fetch\_add(1, SeqCst); assert\_eq!(prev, 10); let prev = number.swap(2, SeqCst); assert\_eq!(prev, 11); assert\_eq!(number.load(SeqCst), 2);

#### std::sync::Mutex

```
let lock = Mutex::new(vec![1, 2, 3]);
{
    let mut vector = lock.lock();
    vector.push(3);
}
// no more access to `vector`,
// lock is unlocked
```

#### crossbeam

- Epoch-based memory reclamation
- Easy translation of algorithms that require GC
- Work stealing deque
- MPMC queues

#### rayon



#### rayon



#### 100% Safe

- Everything you just saw is foolproof
- No segfaults
- No data races
- No double frees...

#### Concurrency?

Rust?

Libraries

**Futures** 

#### Async I/O in Rust (last year)

- *mio,* a "cross platform epoll" event loop library
- Servers were hand-written state machines
- Composition was quite difficult





#### 📮 facebook / wangle

#### What's a future?

- Database query
- RPC request
- Timeouts
- CPU intensive work
- Socket readiness

#### What's a future in Rust?

#### trait Future {

}

type Item;
type Error;

fn poll(&mut self) -> Poll<Item, Error>;
// ...

### Composing futures

// Run one future, then another
f.and\_then(|v| new\_future(v))

// Wait for one of two futures
a.select(b)

// Wait for both futures
a.join(b)

# Async I/O in Rust (today)

- *futures*, a foundational abstraction for Async I/O
- Tokio, a runtime built on *mio* and *futures*
- Futures are **at all layers** of the stack

#### Zero-cost futures

- No allocations in combinators
- No synchronization in combinators
- Library is #![no\_std] compatible
- One dynamic dispatch per event
- One allocation per connection



Reqs/s

### Don't just take my word

Tokio

Framework	Best performance (higher is better)		Cls	Lng	Plt	FE	Aos	IA	Errors
octane	4,287,990	100.0%	Plt	c	Non	Non	Lin	Rea	112
tokio-minihttp	4,040,435	94.2%	Mcr	Rus	Rus	tok	Lin	Rea	
libreactor	3,912,028	91.2%	Plt	c	lib	Non	Lin	Rea	τ
<b>rapidoid-http-fast</b>	3,860,425	90.0%	Pit	Jav	Rap	Non	Lin	Rea	
<b>_</b> rapi <u>do</u> id	3,832,927	89.4%	Plt	Jav	Rap	Non	Lin	Rea	T
undertow	3,241,383	75.6%	Pit	Jav	Utw	Non	Lin	Rea	
colossus	3.173.481	74.0%	Mcr	Sca	Akk	Non	Lin	Rea	D

#### trait Future

- Allows for most specialized implementation
- Enables inter-combinator optimizations
- Permits dynamic dispatch when required

#### trait Future



#### Cancellation

- Cancel a future by dropping it
- Ownership implies one handle to a future
- Deterministic destruction so we know what drops

# Questions?